

Considering Execution Environment Resilience: A White-Box Approach

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Honeywell

What is this all about?

Considering
Execution
Environment
Resilience:
A White-Box
Approach

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Berges,
D. Buchs

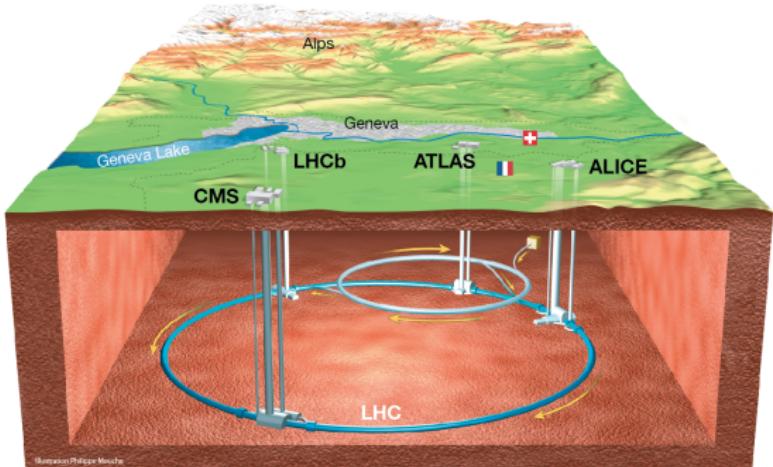
How to

- ▶ generate test cases w. little user interaction
- ▶ on a large scale
- ▶ unit/component level

Welcome to CERN

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Credit: CERN (www.cern.ch)

Welcome to CERN

Considering
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- ▶ LHC, experiments, infrastructure (e.g. power grid)
- ▶ large-scale, widespread, complex systems
- ▶ many types of hard- and software
- ▶ > 100 subsystems,
10,000s of devices,
100,000s of measurement points
- ▶ thousands of physicists/engineers/workers

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high reliability and resilience expectations

How do we supervise it?

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SIMATIC WinCC Open Architecture

- ▶ two frameworks (UNICOS, JCOP) built on top
- ▶ *Control* (CTRL): proprietary scripting language

So where is the problem?

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- ▶ until recently no automated unit test support
- ▶ frequent changes in execution environment
- ▶ (mostly) manual verification
- ▶ big expenses (time) on QA side

Testing

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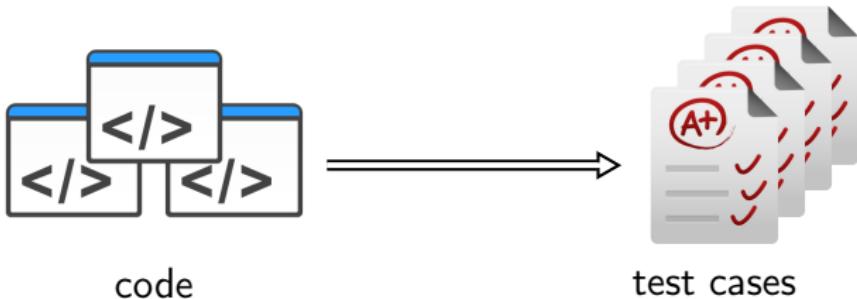
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```
1 f(x){  
2     if GLOBAL_VAR:  
3         return dbGet(x)  
4     else:  
5         return -1  
6 }
```

f(x)

```
1 test_f(){  
2     dbSet("test",5) // prepare  
3     GLOBAL_VAR = True  
4     x = f("test") // act  
5     assert(x == 5) // assert  
6 }
```

Test case for f(x)



Iterative TEst Case System

- ▶ regression testing
- ▶ consider dependencies
- ▶ automatic test case generation (ATCG)
- ▶ build on existing research & tools
- ▶ generate unit & component tests

Automated Test Case Generation

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- ▶ source code based
- ▶ black-box (function signature) vs. white-box (function body)

Semi-purification

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- replace dependencies with parameters

```
1 f(x){  
2     if GLOBAL_VAR:  
3         return dpGet(x)  
4     else:  
5         return -1  
6 }
```

A non-pure function

```
1 f_sp(x,a,b){  
2     if a:  
3         return b  
4     else:  
5         return -1  
6 }
```

Semi-purified $f(x)$

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Semi-purification

- replace dependencies with parameters

```
1 f(x){  
2     if GLOBAL_VAR:  
3         return dpGet(x)  
4     else:  
5         return -1  
6 }
```

A non-pure function

```
1 f_sp(x,a,b){  
2     if a:  
3         return b  
4     else:  
5         return -1  
6 }
```

Semi-purified $f(x)$

```
1 test_f_sp(){  
2     x = f("test",True,5) // act  
3     assert(x == 5) // assert  
4 }
```

Test case

Semi-purification (cont.)

Considering
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- replace dependencies with parameters

```
1 functionA(x){  
2     a = functionB(x)  
3     return a  
4 }  
5  
6 functionB(x){  
7     b = GLOBAL_VAR  
8     b++  
9     return b  
10 }
```

Function with SRC

```
1 functionA_sp(x,y){  
2     a = functionB(x,y)  
3     return a  
4 }  
5  
6 functionB_sp(x,y){  
7     b = y  
8     b++  
9     return b  
10 }
```

Semi-purified w. SRC

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Semi-purification (cont.)

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- replace dependencies with parameters

```
1 functionA(x){  
2     a = functionB(x)  
3     return a  
4 }  
5  
6 functionB(x){  
7     b = GLOBAL_VAR  
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```

Function with SRC

```
1 functionA_sp(x,y){  
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4 }  
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6 functionB_sp(x,y){  
7     b = y  
8     b++  
9     return b  
10 }
```

Semi-purified w. SRC

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Semi-purification: Concept

- ▶ code contains dependencies
 - ▶ global variables, data base values, subroutine calls, other resources
- ▶ manual way: test doubles (mocks, stubs, fakes, . . .) [ME06]
- ▶ remove dependencies
 - ▶ based on localization [SW03, SK13]
 - ▶ input parameters instead of dependencies
 - ▶ use any ATCG (black- and white-box)

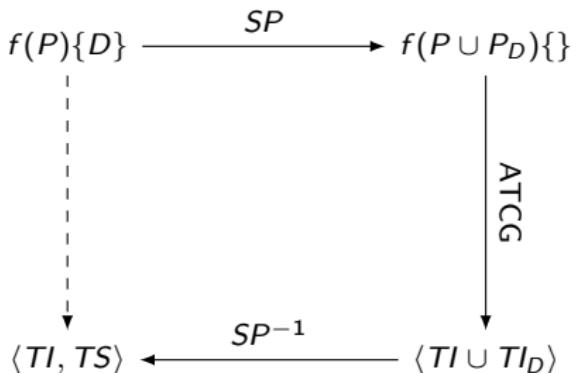


Figure: Test case generation schema

- ▶ P ... Parameters
- ▶ D ... Dependencies

- ▶ TI ... Test Input
- ▶ TS ... Test Setup Routine

Identified Bottlenecks

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- ▶ Loops
- ▶ Shared Subroutines
- ▶ Concurrency

SP: Bottlenecks (cont.)

► Loops

```
1 sleepUntilReady(){ // a = bool
2
3     while dpGet(notReadyDP):
4         sleep(5) // sleep for 5 seconds
5
6 }
```

A semi-purified loop

SP: Bottlenecks (cont.)

► Loops

```
1 sleepUntilReady(a){ // a = bool
2
3     while a:// replaces dpGet(notReadyDP)
4         sleep(5) // sleep for 5 seconds
5
6 }
```

A semi-purified loop

Test Cases:

- a: False \Rightarrow loop not executed
- a: True \Rightarrow endless loop

SP: Bottlenecks (cont.)

► Loops

```
1 sleepUntilReady(a){ // a = [bool]
2     i = 0
3     while a[i]:// replaces dpGet(notReadyDP)
4         sleep(5) // sleep for 5 seconds
5         i++
6 }
```

A semi-purified loop

Test Cases:

- a: [False] ⇒ loop not executed
- a: [True, True, . . . , False] ⇒ loop execution

Questions:

- how long should the list be?
- how to modify correctly?
- Test modified code or w. threads?

SP: Bottlenecks (cont.)

► Shared subroutine dependencies

```
1 var SPEED_VAR = 1
2 adjustSpeed(){
3     x = getTheSpeed()
4     if x < 10 :
5         doubleTheSpeed()
6 }
```

CUT

```
1 getTheSpeed(){
2     return SPEED_VAR
3 }
```

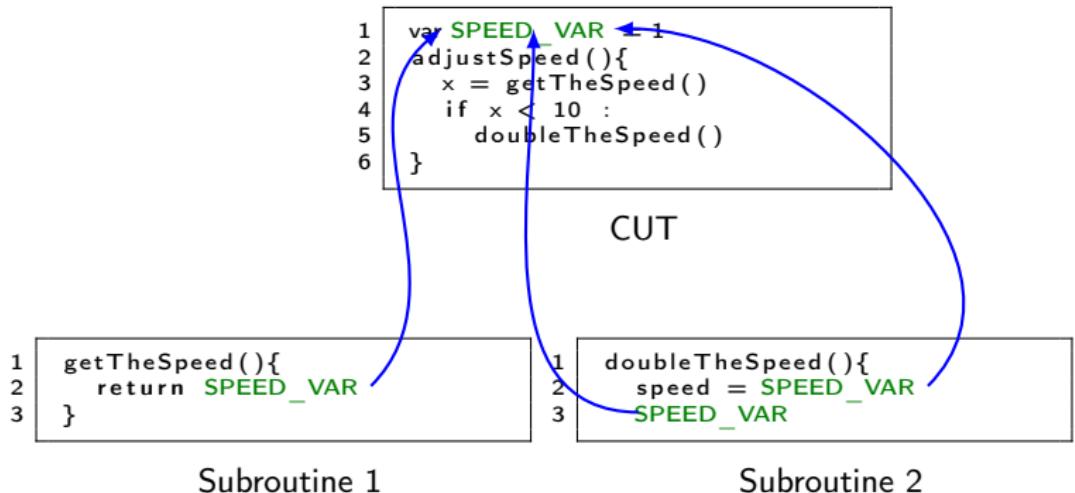
Subroutine 1

```
1 doubleTheSpeed(){
2     speed = SPEED_VAR
3     SPEED_VAR
```

Subroutine 2

SP: Bottlenecks (cont.)

- ▶ Shared subroutine dependencies



SP: Bottlenecks (cont.)

► Shared subroutine dependencies

```
1  adjustSpeed(a){  
2      x = getTheSpeed(a)  
3      if x < 10 :  
4          doubleTheSpeed()  
5      }  
6  }
```

CUT

```
1  getTheSpeed(a){  
2      return a // SPEED_VAR  
3  }
```

Subroutine 1

```
1  doubleTheSpeed(){  
2      speed = SPEED_VAR  
3      SPEED_VAR
```

Subroutine 2

SP: Bottlenecks (cont.)

► Shared subroutine dependencies

```
1  adjustSpeed(a, b){  
2      x = getTheSpeed(a)  
3      if x < 10 :  
4          doubleTheSpeed(b)  
5      }  
6  }
```

CUT

```
1  getTheSpeed(a){  
2      return a // SPEED_VAR  
3  }
```

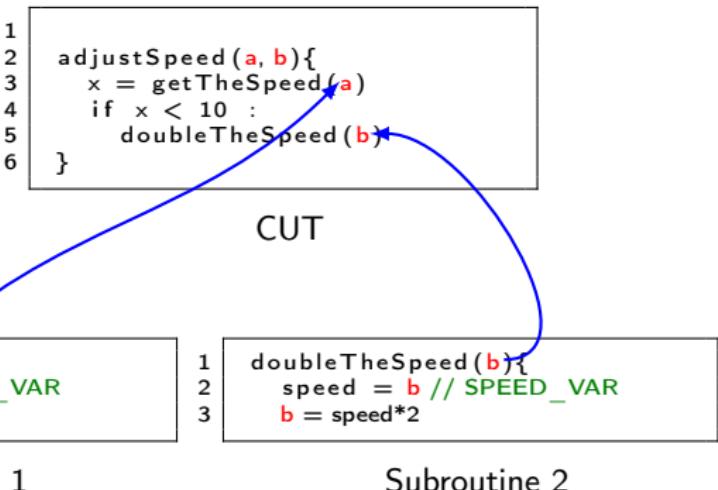
Subroutine 1

```
1  doubleTheSpeed(b){  
2      speed = b // SPEED_VAR  
3      b = speed*2
```

Subroutine 2

SP: Bottlenecks (cont.)

- ▶ Shared subroutine dependencies



SP: Bottlenecks (cont.)

- ▶ Shared subroutine dependencies

```
1  adjustSpeed (a,b){  
2      x = getTheSpeed (a)  
3      if x < 10 :  
4          doubleTheSpeed (b)  
5      }  
6 }
```

CUT

```
1  getTheSpeed (a){  
2      return a // SPEED_VAR  
3 }
```

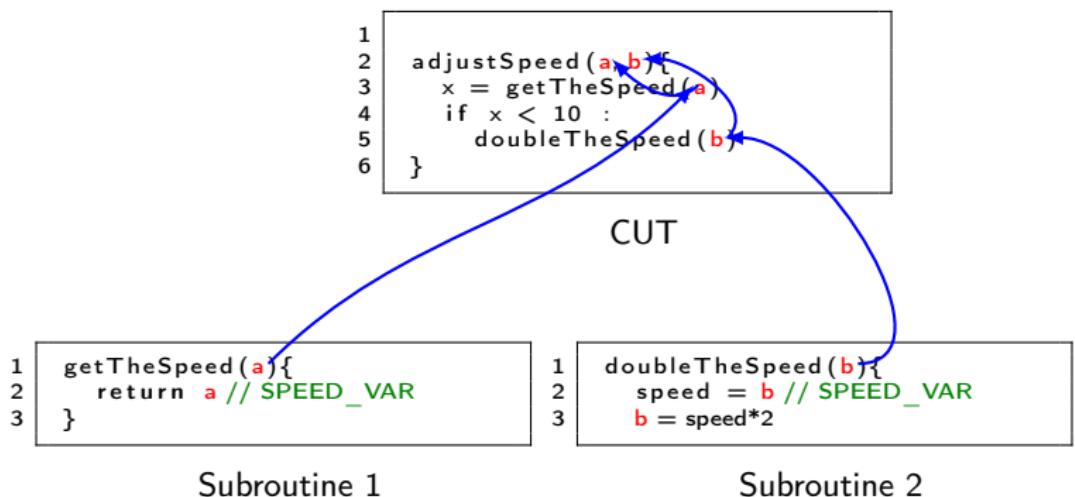
Subroutine 1

```
1  doubleTheSpeed (b){  
2      speed = b // SPEED_VAR  
3      b = speed*2
```

Subroutine 2

SP: Bottlenecks (cont.)

- ▶ Shared subroutine dependencies



- ▶ BUT: **a** and **b** replace the same dependency

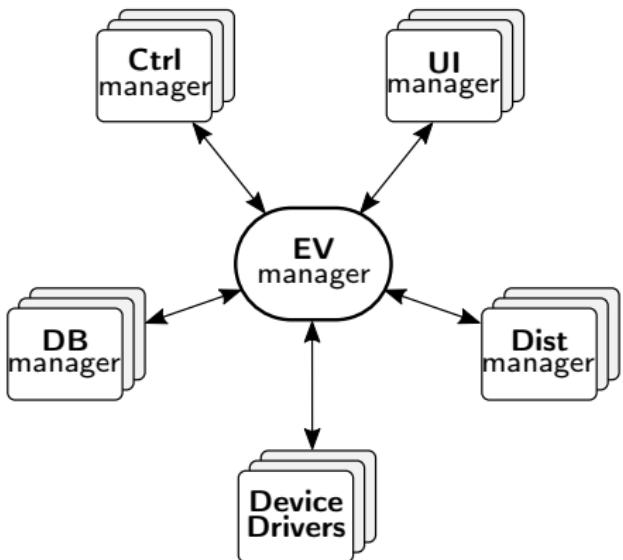


Figure: WinCC OA's manager concept

SP: Bottlenecks (cont.)

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- ▶ access to ALL data points
- ▶ 100+ sub-systems
- ▶ discover dirty read/write scenarios (`adjustSpeed()`)

ITEC implementation

ITEC

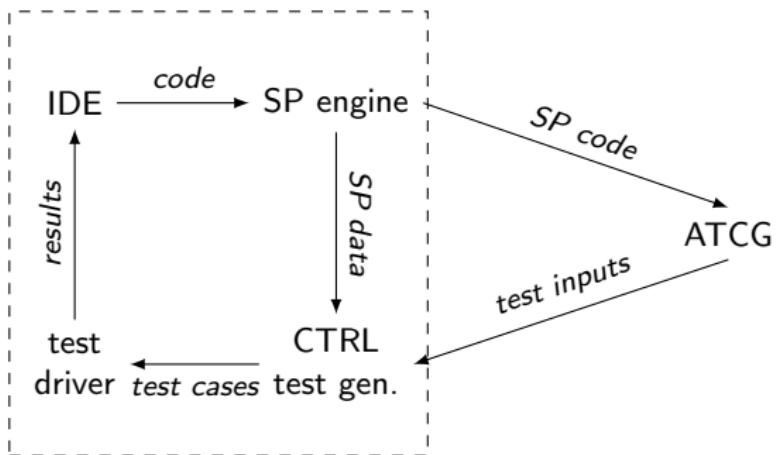


Figure: ITEC workflow

ITEC implementation

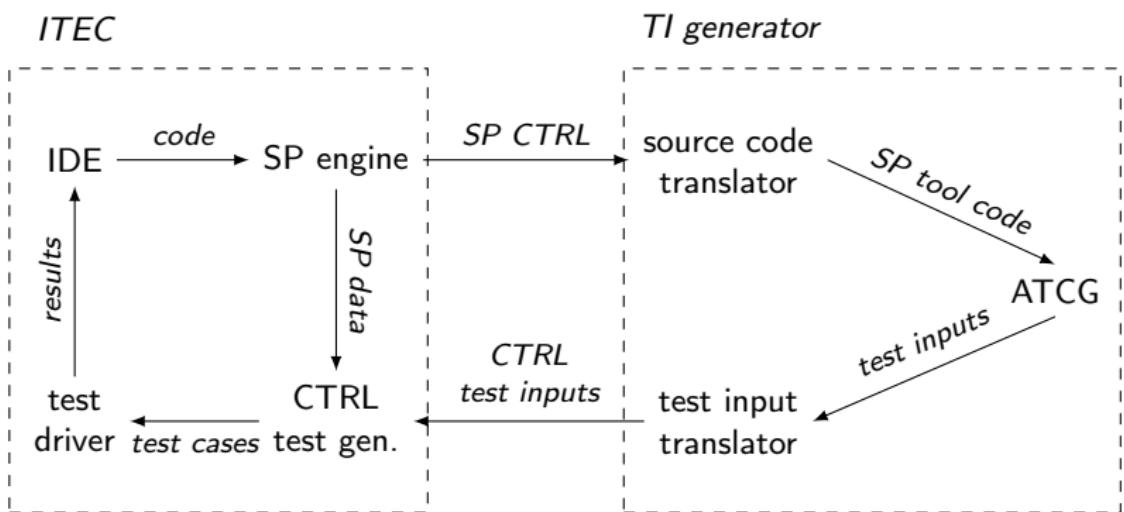


Figure: ITEC workflow

Conclusion

Semi-purification

- ▶ remove dependencies
- ▶ facilitate unit tests generation
- ▶ use black-box techniques on all code

Future works

What's next?

- ▶ connect to test framework & roll-out
- ▶ intrinsic domain information
- ▶ compare different ATCG
- ▶ research into bottlenecks

References

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